Secure BFD Sequence Numbers
draft-ietf-bfd-secure-sequence-numbers-01

Abstract

This document describes a security enhancements for the BFD packet’s sequence number.

Requirements Language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

Status of This Memo

This Internet-Draft is submitted in full conformance with the provisions of BCP 78 and BCP 79.

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1. Introduction

BFD [RFC5880] section 6.7 describes the use of monotonically incrementing 32-bit sequence numbers for use in authentication of BFD packets. While this method protects against simple replay attacks, the monotonically incrementing sequence numbers are predictable and vulnerable to more complex attack vectors. This document proposes the use of non-monotonically-incrementing sequence numbers in BFD authentication TLVs to enhance the security of BFD sessions. Specifically, the document presents a method to generate pseudo-random sequence numbers on the frame by algorithmically hashing monotonically increasing sequence numbers. Further security may be introduced by resetting un-encrypted sequence to a random value when the 32-bit sequence number rolls-over.

2. Theory of operations

Instead of monotonically increasing the sequence number or even occasionally monotonically increasing the sequence number, the next sequence number is generated by computing a hash on what would have been the next sequence number using a shared key. That computed hash is then inserted into the sequence number field of the packet. In case of BFD Authentication [I-D.ietf-bfd-optimizing-authentication], the sequence number used in computing an authenticated packet would be this new computed hash. Even though the BFD Authentication
sequence number is independent of this enhancement, it would benefit by using the computed hash.

A normal BFD packet with authentication will undergo the following steps, where:

- **[O]**: original RFC 5880 packet with monotonically increasing sequence number
- **[S]**: pseudo random sequence number
- **[A]**: Authentication

<table>
<thead>
<tr>
<th>Sender</th>
<th>Receiver</th>
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<tbody>
<tr>
<td>[O] [S] [A]</td>
<td>[A] [S] [O]</td>
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</table>

In order to encode a sequence number, the sender would identify a hash algorithm (symmetric) that would create a 32 bit hash. The hashing key is provisioned securely on the sender and receiver of the BFD session. The mechanism of provisioning such a key is outside the scope of this draft. Instead of using the sequence number, the sender encodes the sequence number with the hashing key to produce a hash.

Upon receiving the BFD Control packet, the receiver compares the received sequence number against the expected sequence number. The mechanism used for comparing is an implementation detail (implementations may pre-calculate the expected hashed sequence number, or decrypt the received sequence number before comparing against expected value). To tolerate dropped frames, the receiver MUST compare the received sequence number against the current expected sequence number (previous received sequence number + 1) and N subsequent expected sequence numbers (where N is greater than or equal to the detect multiplier). Note: The first sequence number can be obtained using the same logic as the My Discriminator value.

- **k**: hashing key
- **s**: sequence number
- **O**: original RFC 5880 packet with monotonically increasing sequence number
- **R**: remainder of packet
- **H1**: hash of s
H2: hash of entire packet
A: H2 + insertion in packet
hash(s, k) = H1
hash((H1 + R), k) = H2
hash’(H1, k) == s ? Good sequence number : bad sequence number

Sender                Receiver
[O] [H1] [A] -------- [A] [H1] [O]

3. Impact of using a hash

Under this proposal, every packet’s sequence number is encoded within a hash. Therefore there is some impact on the system and its performance while encoding/decoding the hash. As security measures go, this enhancement greatly increases the security of the packet with or without authentication of the entire packet.

4. IANA Considerations

This document makes no request of IANA.

Note to RFC Editor: this section may be removed on publication as an RFC.

5. Security Considerations

6. Acknowledgements

7. References

7.1. Normative References


7.2. Informative References

[I-D.ietf-bfd-optimizing-authentication]
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