Yang for I2RS Protocol
draft-hares-netmod-i2rs-yang-00.txt

Abstract

This document requests one yang model addition that will support ephemeral state and provides notes for the implementers who wish to implement ephemeral state for the I2RS Protocol. The purpose of this document is to provide implementers of ephemeral state with background and open issues that they should consider when implementing ephemeral state that satisfies the I2RS protocol.

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This a proposal for yang additions to support the first version of the I2RS protocol.

The I2RS architecture [RFC7921] defines the I2RS interface "a programmatic interface for state transfer in and out of the Internet routing system". The I2RS protocol is a protocol designed to a higher level protocol comprised of a set of existing protocols which have been extended to work together to support a new interface to the routing system. The I2RS protocol is a "reuse" management protocol which creates new management protocols by reusing existing protocols and extending these protocols for new uses, and has been designed to be implemented in phases [RFC7921].
The first version of the I2RS protocol is comprised of extensions to existing features of NETCONF [RFC6241] and RESTCONF [I-D.ietf-netconf-restconf]. The data modeling language for the I2RS protocol will be Yang [RFC7950] with features and extensions proposed in this draft.

The structure of this document is:

Section 2 provides definitions for terms in this document.

Section 3 summarizes the changes to configuration data store, NETCONF, RESTCONF, and YANG.

[I-D.ietf-i2rs-ephemeral-state] specifies the I2RS requirements for the ephemeral state. Section 4 discusses how these requirements might be implemented in a control plane datastore.

Section 5 describes the one required Yang model addition for I2RS (ephemeral key word). This section also describes elements of information in the NETCONF/RESTCONF implementations that must be queryable by the I2RS protocol implementations.

2. Definitions Related to Ephemeral Configuration

This section reviews definitions from I2RS architecture [RFC7921] and NETCONF operational state definitions [I-D.nmdsdt-netmod-revised-datastores] before using these to construct a definition of the ephemeral data store.

2.1. Requirements language

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in RFC 2119 [RFC2119].

2.2. I2RS Definitions

The I2RS architecture [RFC7921] defines the following terms:

ephemeral data: is data which does not persist across a reboot (software or hardware) or a power on/off condition. Ephemeral data can be configured data or data recorded from operations of the router. Ephemeral configuration data also has the property that a system cannot roll back to a previous ephemeral configuration state. (See [RFC7921] for an architectural overview, [I-D.ietf-i2rs-ephemeral-state] for requirements, and [I-D.nmdsdt-netmod-revised-datastores] for discussion of how the ephemeral datastore as a control plane datastore interacts with
intended datastore and dynamic configuration protocols to form the applied datastore”.

local configuration: is the data on a routing system which does persist across a reboot (software or hardware) and a power on/off condition. Local configuration has the ability to roll back to a previous configuration state. Local configuration is defined as the intended datastore [I-D.nmdsdt-netmod-revised-datastores] which is modified by dynamic configuration protocols (such as DHCP) and the I2RS ephemeral data store.

dynamic configuration protocols datastore are configuration protocols such as DHCP that interact with the intended datastore (which does persist across a reboot (software or hardware) power on/off condition), and the I2RS ephemeral state control plane datastore.

operator-applied policy: is a policy that an operator sets that determines how the ephemeral datastore as a control plane datastore interacts with applied datastore (as defined in [I-D.nmdsdt-netmod-revised-datastores]). This operator policy consists of policy knobs that the operator sets to determine how the I2RS agent control plane ephemeral state datastore will interact with the intended configuration datastore and the dynamic configuration protocol datastore. Three policy knobs could be used to implement this policy:

* policy knob 1: I2RS Ephemeral control-plane datastore takes precedence over the intended datastore in the routing protocols.

* policy knob 2: Updated intended configuration datastore takes precedence over the I2RS ephemeral control-plane data store in the routing protocols

* policy knob 3: Ephemeral control plane datastore takes precedence over any other dynamic configuration protocols datastore.

An practical example for three states of the operator-applied policy may help the reader understand the concept. Consider the following three desired outcomes with their policy knob states:

Monitoring Features only  The policy knob settings are:

Policy knob 1=false,

policy knob 2=true,
Policy knob 3=false,

Action: I2RS protocol software feature is installed, but the operator does not want the I2RS ephemeral datastore to take precedence (that is be used) on any variables in the applied configuration datastore. This policy set might be valid if I2RS is only suppose to monitor data on this node through newly defined parameters.

I2RS Agent Changes win the policy knob settings would be:

Policy knob 1=true,

policy knob 2=false,

Policy knob 3=false,

Action: This is the normal case for the I2RS Agent where the ephemeral control-plane datastore takes precedence over the intended configuration datastore and dynamic configuration datastores. The values from the I2RS ephemeral datastore are used rather than the intended configuration datastore and the dynamic configuration protocol datastore. When the ephemeral data is removed by the I2RS agent, the dynamic configuration datastore and the intended configuration datastore state is restored, combined and passed to the routing protocols for application.

Just change until next configuration update the policy knob settings would be:

Policy knob 1=true,

policy knob 2=true,

Policy knob 3=false,

Action: This case can occur if the I2RS Client write to the ephemeral control plane data store is only suppose to take precedence until the next configuration cycle from a centralized system. Suppose the local configuration is get by the centralized system at 11:00pm each night. The I2RS Client writes temporary changes to the routing system via the I2RS agent ephemeral write. At 11:00pm, the local configuration update overwrite the ephemeral. The I2RS Agent notifies the I2RS Client which is tracking which of the ephemeral changes are being overwritten.
3. Overview of Changes

This overview reviews the following:

- What NETCONF [RFC6241] protocol existing features required for I2RS protocol and what extension for these extension features that are needed for the I2RS protocol version 1.

- What RESTCONF [I-D.ietf-netconf-restconf] protocol existing features are required for the I2RS protocol and what extensions are needed for I2RS protocol version 1.

- An overview of the Yang 1.1 data modeling language [RFC7950] features are needed for I2RS protocol version 1.

- An overview of the extensions to Yang 1.1 data modeling language [RFC7950] that are needed for the I2RS protocol version 1.

3.1. I2RS protocol requirements

The requirements for the I2RS protocol are defined in the following documents:

- I2RS Problem Statement [RFC7920],
- I2RS Architecture [RFC7921],
- I2RS Traceability [RFC7922],
- Publication and Subscription [RFC7923],
- I2RS Ephemeral State Requirements, [I-D.ietf-i2rs-ephemeral-state]
- I2RS Protocol Security Requirements, [I-D.ietf-i2rs-protocol-security-requirements]

The Interface to the routing System (I2RS) creates a new capability for the routing systems, and with greater capabilities come a greater need for security. The requirements for a secure environment for I2RS is described in [I-D.ietf-i2rs-security-environment-reqs].

3.2. NETCONF Features and Extensions

The features the I2RS protocol requires are:

- NETCONF [RFC6241] with its updates [RFC7803],
o Network Access Control Model [RFC6536] with update (draft-bierman-netconf-rf6536bis)

o Running NETCONF over TLS with mutually X.509 authentication [RFC7589]

o Keystore Model [I-D.ietf-netconf-keystore],

o Subscribing to Yang Datastore updates [I-D.ietf-netconf-yang-push],

o NETCONF support for Event Notifications [I-D.ietf-netconf-netconf-event-notifications],

o Subscribing to NETCONF Events (updated) [I-D.ietf-netconf-rfc5277bis]

o Yang Patch Media type [I-D.ietf-netconf-yang-patch],

o NETCONF/RESTCONF Zero Touch provisioning [I-D.ietf-netconf-zerotouch],

o TLS Client and Server Models [I-D.ietf-netconf-restconf-client-server]

o Call Home [I-D.ietf-netconf-call-home],

o Module library [RFC7895],

o NETCONF/RESTCONF Zero Touch provisioning [I-D.ietf-netconf-zerotouch],

3.3. RESTCONF features and Extensions

This protocol strawman utilizes the following existing proposed features for NETCONF and RESTCONF

o RESTCONF [I-D.ietf-netconf-restconf]

o Module library [RFC7895],

o Publication/Subscription via Push [I-D.ietf-netconf-yang-push],

o Patch [I-D.ietf-netconf-yang-patch],

o syslog yang module (both [RFC5424] and [I-D.ietf-netmod-syslog-model]
3.4. Assumptions on Data Store Model Melee

The NETMOD Working Group has been working to create new definitions of datastores based on feedback from operators on desiring a split between operational state and configuration state.

This document takes [I-D.nmndsdt-netmod-revised-datastores] as the current status of the datastore discussion on configuration state, operational state, ephemeral state changes (via I2RS), and routing protocol state. The following things need to be carefully defined in this work:

What is a dynamic configuration protocol (is it I2RS or DHCP)

What is a control-plane datastore - (ephemeral state only or others? )

How to express the policy knobs that provide preference between intended configuration, control plane datastore, and dynamic configuration protocols

How does operational state allow for operational state to be defined by ephemeral-only data models, and mixed (ephemeral + intended configuration)

[I-D.nmndsdt-netmod-revised-datastores] is making good progress, but these additional details need to be tied down.

4. Ephemeral Data

This section provides an overview of the ephemeral data store as a control plane datastore and discusses several concepts that implementers need to consider and provide feedback on. The concepts include basic ephemeral datastore concepts, I2RS caching of ephemeral data, issues for massive data flow, error handling (normal and reduced), use of IPFIX or Binary for carrying I2RS ephemeral data, and ephemeral state.

This section augments [I-D.nmndsdt-netmod-revised-datastores] to begin to discuss how the ephemeral state control-plane datastore might be implemented.

The purpose of this section is to gather implementer wisdom on the ephemeral datastore into one place. This section discusses:

Ephemeral state as a control plane data store

Qualities of ephemeral datastores
Need to support Massive amounts of configuration data,
Two types of Error handling (regular, reduced)
Should we support link to IPFIX in I2RS protocol and ephemeral state?
Binary encoding for RESTCONF/NETCONF
Ephemeral state in DDoS environments.

[I-D.ietf-i2rs-ephemeral-state] describes the requirements for I2RS ephemeral state.

This section augments [I-D.nmdsdt-netmod-revised-datastores] to begin to discuss how the ephemeral state control-plane datastore might be implemented. This initial draft refines the general description so that early I2RS ephemeral state implementations may progress.

4.1. Ephemeral Control Plane Datastore

[I-D.nmdsdt-netmod-revised-datastores] architecture suggests that the applied configuration is the combination of intended datastore, the dynamic configuration protocols, and the control-plane datastores. As described above, there are policy knobs which allow the I2RS Agent to handle deciding what specific configuration variables is installed in protocols (e.g. BGP) or protocol independent functions (RIB or Filters). In addition, the control-plane datastore may store the parameters need to provide publication of events, statistics, telemetry within the ephemeral control-plane datastore.

The ephemeral data-store may have models which learn operational state and augment it by configuration. For example [I-D.ietf-i2rs-yang-l3-topology] uploads ospf and isis topology information from the routing system and allows configuration of additional links or nodes.

This new architecture is a multiple panes-of-glass model where the decision on what value is chosen is based on policy. The extension of this model is that it is possible for two or more of the control-plane datastores to be ephemeral. If this occurs, then the policy knobs must define the how the 2+ ephemeral datastores interact with each other and the configuration state.
4.2. Qualities of Ephemeral Datastore

Note: The requirements for ephemeral state are in: [I-D.ietf-i2rs-ephemeral-state]).

This section provides a discussion so that implementers writing code for these datastores can discuss what needs to be standardized and what does not need to be standardized.

The ephemeral data store has the following general qualities:

1. Ephemeral state is not unique to I2RS work.
2. The ephemeral datastore is never locked.
3. The ephemeral portion of the intended configuration, applied state, and derived state does not persist over a reboot,
4. an ephemeral node cannot roll-back to its previous value,
5. Since ephemeral data store is just data that does not presist over a reboot, then in theory any node or group of nodes in a YANG data model could be ephemeral. The YANG data module must indicate what portion of the data model (if any) is ephemeral.
   * A YANG data module could be all ephemeral (e.g. [I-D.ietf-i2rs-rib-data-model]) with no directly associated configuration models,
   * A YANG model could be all ephemeral but associated with a configuration model
   * or a single data node or data tree could be made ephemeral.
6. The management protocol (NETCONF/RESTCONF) needs to signal which portions of a data model(node, tree, or data model) are ephemeral in the module library [RFC7895].

4.3. I2RS Agent Caching of Ephemeral Data

The multiple control-plane datastore model [I-D.nmdsdt-netmod-revised-datastores] architecture allows multiple datastores which could allow an implementation of caching of ephemeral data in the I2RS Agent by having a main and a backup I2RS agent. Early implementations should at least support the single ephemeral data model, but MAY support the multiple datastore mode. It is important that these early implementations provide feedback for standardization on the following:
the policy knobs needed to make single ephemeral control planes datastores function,

the policy knobs need to make multiple ephemeral control plane datastores which support caching work.

4.4. Massive Amounts of Configuration Data

Large amounts of data can flow from the I2RS agent to the I2RS client, or from the I2RS client to the I2RS Agent. The I2RS client may set or query ephemeral configuration in the routing system via the I2RS agent and receive operational state, notifications, or logging from the I2RS Agent on behalf of the I2RS routing system. I2RS Clients can send large amount of ephemeral configuration data to the I2RS Agent. The writes may be done via NETCONF (<edit-config> or an rpc function), or via RESTCONF (PUT, PATCH, POST). Reads can be done via NETCONF <get-config> or RESTCONF GET or query.

The I2RS RIB Data Model [I-D.ietf-i2rs-rib-data-model] also supports the use of rpc to add/delete RIBs, add/delete/update routes, and add/delete nexthops. If the I2RS client does a small to medium number of writes to the I2RS ephemeral state in the I2RS Agent in a routing system, the full validation that NETCONF or RESTCONF does will be able to be done without any reduction in speed to the I2RS high-performance system. For example, if the I2RS RIB Data Model has adds a 1000 routes, the I2RS RIB use of rpc to add/delete/update routes should be able to provide a high-performance system. Alternatively the NETCONF <edit-config> could update these 1000 routes with a write, or the RESTCONF POST, PUT or PATCH should be able to add the 1000 routes.

If a large number of ephemeral routes or filters are written (updates or new) by the I2RS Client to the ephemeral state in the I2RS agent, one of the key issues for a high performance interface is the time it takes to validate routes. Due to this concern, the I2RS architecture was design to allow less than the full NETCONF or RESTCONF validation. The concept is that the I2RS routes would be validated within the I2RS client and sent via a 99.999% reliable connection. In this scenario, the I2RS Agent would trust the validation that the I2RS Client did, and the communication of the route additions via the network connection.

An experiment regarding this has been done with the ODL code base update of ephemeral routes, but additional experimentation needs to be done prior to finalizing this design. Section 3.4.2 reviews how this process might be done, but many open issues exist in implementing this "low-validation" interface. Without additional experimentation and prototype code, this type of "low-validation"
4.5. Write Error handling

This section reviews I2RS normal error handling and error handling for rpc with no validation checks.

4.5.1. Normal validation checks

An I2RS agent validates an I2RS client’s information by examining the following:

- message syntax validation,
- syntax validation for nodes of data model,
- referential checks (leafref checks MUST clauses, and instance identifier),
- checks groups of data within a data model or groups of data across data models,
- write access to data,
- if write access and values already exist, if I2RS client write access is higher than existing priority.

4.5.1.1. Reduced Validation (Experimental)

Can the I2RS protocol allow for reduced error checking? The need for speed in the I2RS protocol insertions into the I2RS RIB suggest that it is worth experimenting for reduced validation in order to obtain high levels of throughput. If NETCONF or RESTCONF streams pre-checked routes to the datastore, what happens? Implementation experience is needed to determine the feasibility of this approach.

This feature may require a operator-applied policy knob with a "no validation" feature

- operator-applied policy knob enabling this feature;
- rpc in a data model with the yang "ephemeral-validation no-check;"

4.6. IPFIX for traffic monitoring

Due to the potentially large data flow the traffic measurement statistics generate, these statistics are best handled by publication techniques within NETCONF or a separate protocol such as IPFIX. In the future version of the I2RS protocol may desire to support a data
stream outbound from the I2RS Agent to an I2RS client via the IPFIX protocol.

4.7. Binary encoding of RESTCONF/NETCONF

The binary encoding of JSON or XML encoding in RESTCONF or NETCONF may provide a better throughput. Research needs to be done on what is the appropriate binary encoding.

4.8. Ephemeral state in DDoS environments

I2RS ephemeral state may operate in places where there is a DDoS attacks where the network devices are attacked. Is one attack plane the ability to remove all tracing if the I2RS reboots an attack vector?

5. Yang Changes

The data modules supporting the ephemeral datastore can use the Yang module library to describe their datastore.

The following key word must be able to specify ephemeral

   ephemeral true;

Nice to have features:

It would be helpful for implementation of I2RS ephemeral data models to determine if the I2RS protocol feature set can support the I2RS data model needs. For this reason, it is helpful to group protocol features into "versions" and to put flags in the data model. At this point, the best place to put the summary of features is in an data model which defines these features. The discussion between implementers should be whether it is useful to have this features in some general yang location. An example of features that might be needed are:

- i2rs version indicator;
- i2rs transport-nonsecure "ok-to-use";
- i2rs ephemeral-validation nocheck;
- I2rs caching
6. IANA Considerations

This is a protocol strawman - nothing is going to IANA.

7. Security Considerations

The security requirements for the I2RS protocol are covered in [I-D.ietf-i2rs-protocol-security-requirements]. The security environment the I2RS protocol is covered in [I-D.ietf-i2rs-security-environment-reqs]. Any person implementing or deploying the I2RS protocol should consider both security requirements.

8. Acknowledgements

This document is an attempt to distill lengthy conversations on the I2RS proto design team from August

Here’s the list of the I2RS protocol design team members

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